

Formal Security Analysis

Dr. David von Oheimb

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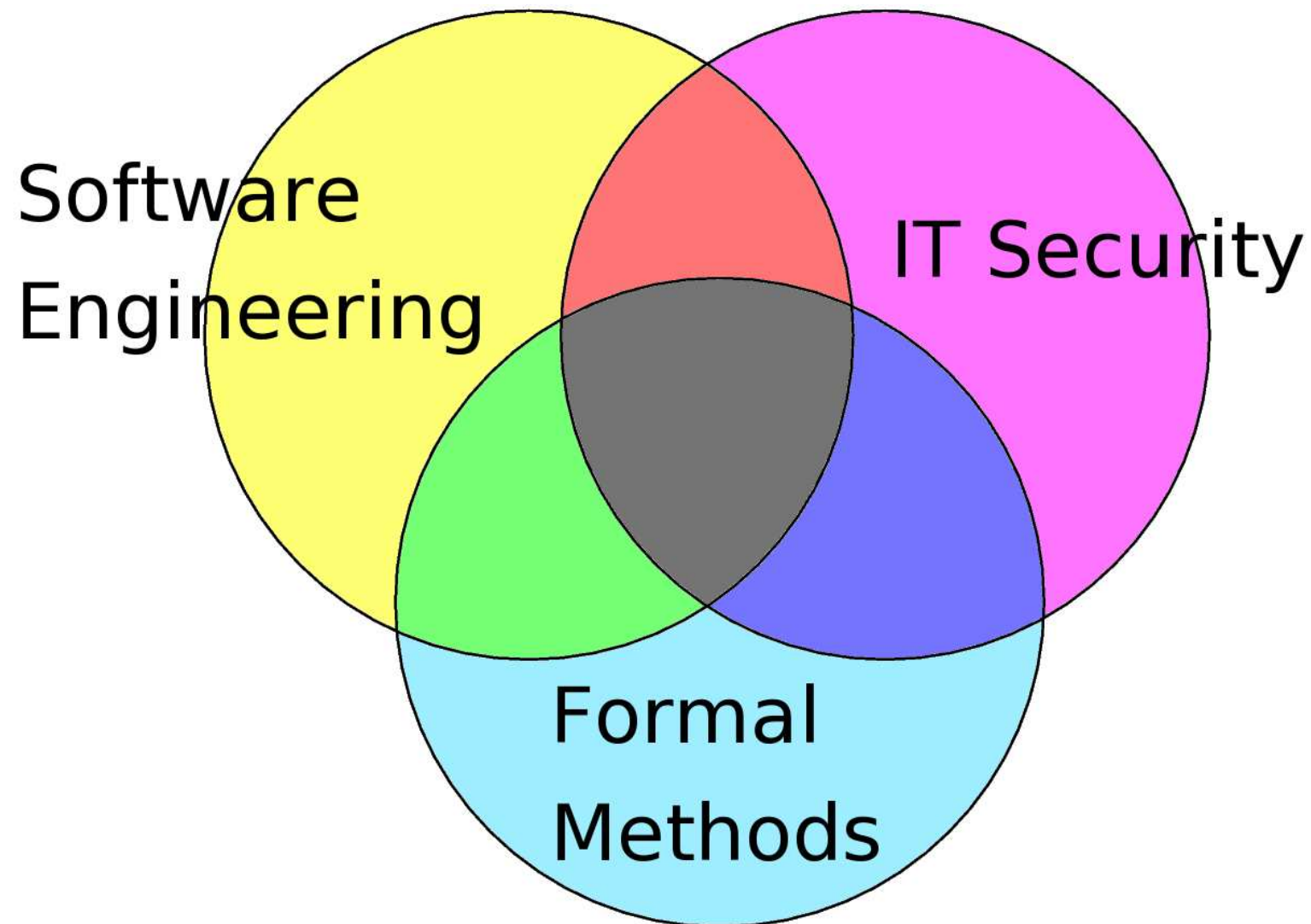


Overview

- What is Formal Security Analysis?
 - IT Security
 - Formal Modeling
 - Practical Considerations
 - Relation to Common Criteria



Formal Security Analysis: Areas



Overview

- What is Formal Security Analysis?

-  **IT Security**

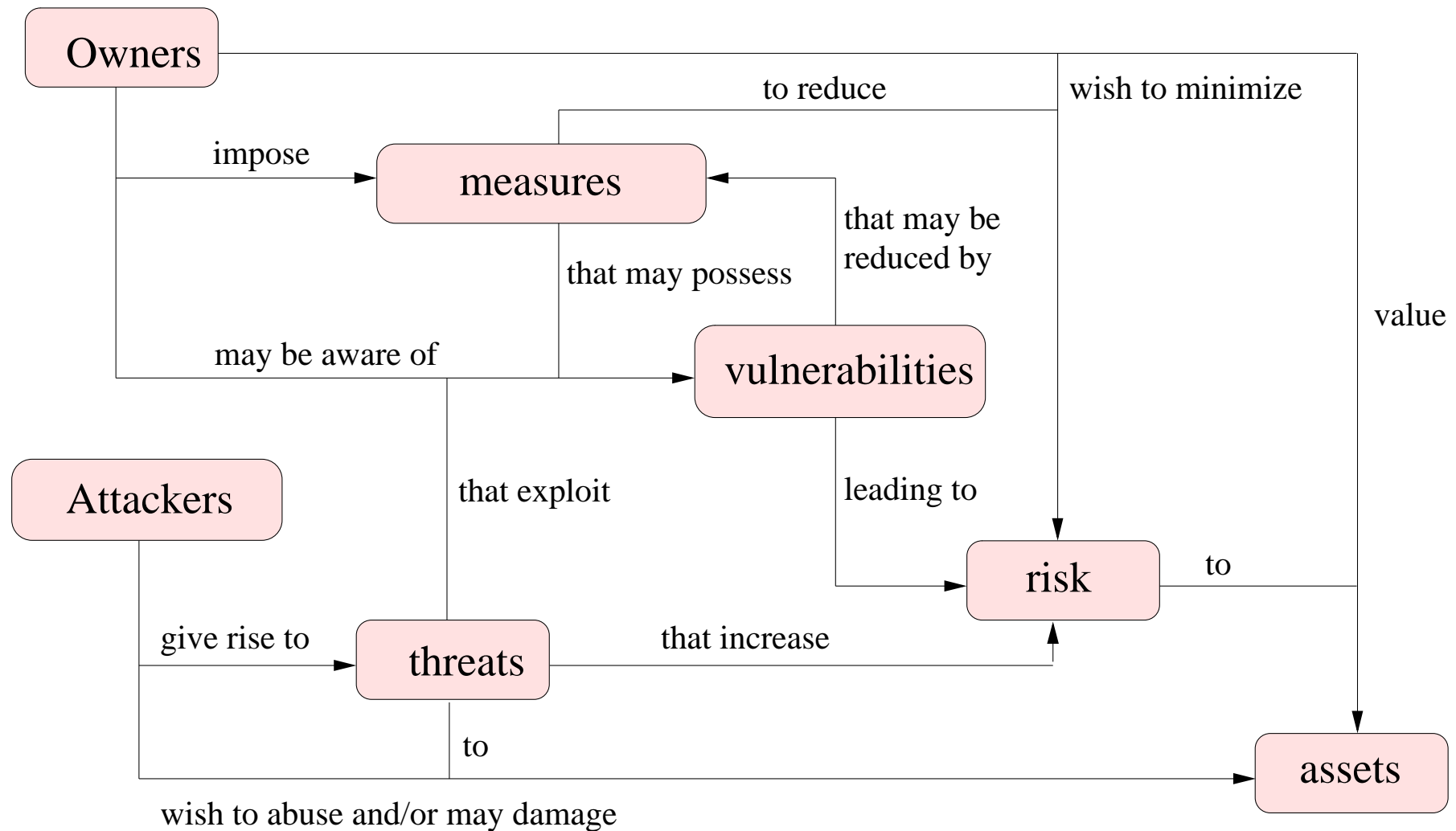
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IT Security

- **IT/Computer security** deals with the prevention, or at least detection, of unauthorized actions by users of a computer system.
 - **Authorization** is central to definition.
 - Sensible only relative to a **security policy**, stating who (or what) may perform which actions.
- Complements **safety**: prevent damage through errors or malfunction



Security Concepts and Relationships



Policy (here implicit) defines authorized actions on assets, i.e., what constitutes legal use (or abuse/damage, respectively).



Security as a Software Engineering Problem

Situation: security loopholes in IT systems will be actively exploited

— in this sense even worse than safety problems!

Goal: achieve absence of attacks by absence of vulnerabilities

— and convince yourself/contractors/customers of this!

Problem: IT systems are very complex, security flaws hard to find.

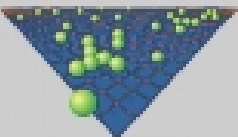
Security cannot be added on, but must be co-designed with the system.

Remedy: address security in all development phases.

Reviews supported by formal security modeling/analysis.

During ...

- **requirements analysis:** helps understanding the security issues
- **design, documentation:** helps improving the quality of specifications
- **implementation:** acts as correctness reference for testing/verification



Goals, Threats, and Mechanisms

Standard breakdown in **security engineering**:

Goals/Objectives: What to achieve

Threats: Which attacks to counter

Mechanisms How to achieve goals

Required for **certification** according to e.g. ITSEC and Common Criteria



Security Goals

- Goals: **CIA**

Confidentiality No unauthorized disclosure/leakage of information

Integrity: No unauthorized modification of information

Availability: No unauthorized impairment of functionality

All these require **authorization** = **authentication** + **access control**.

- Other goals

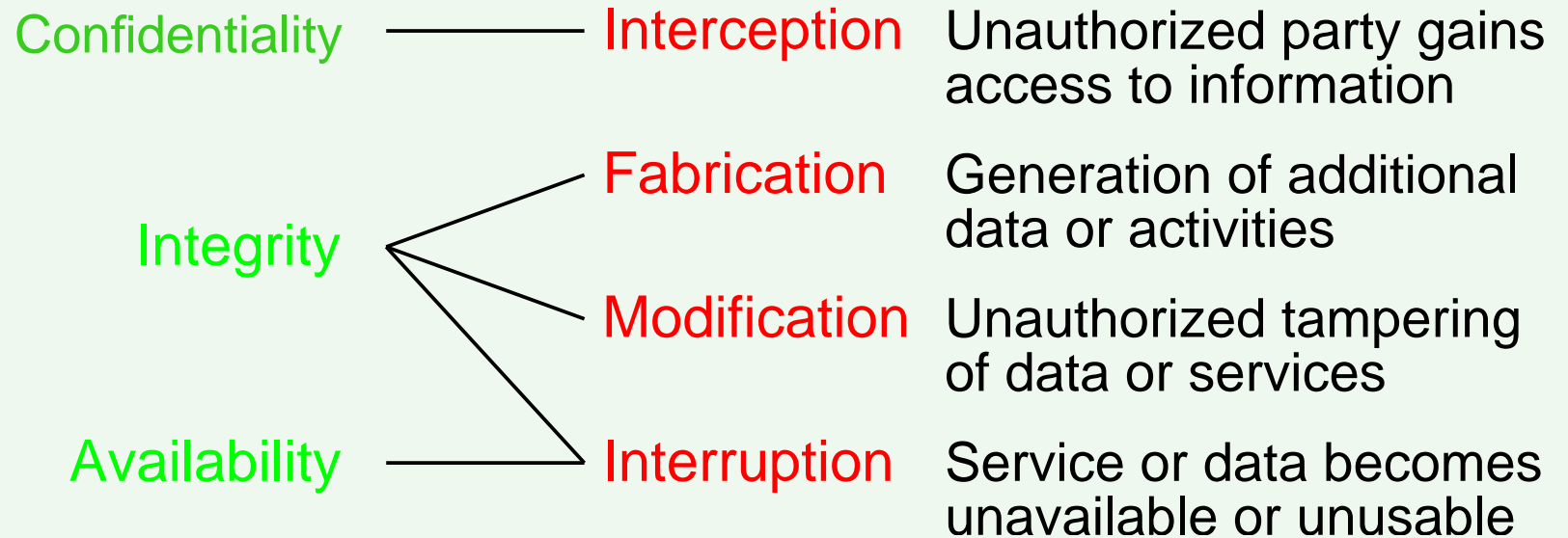
Privacy: User data is only exposed in permitted ways.

Nonrepudiation: One cannot deny responsibility for actions.

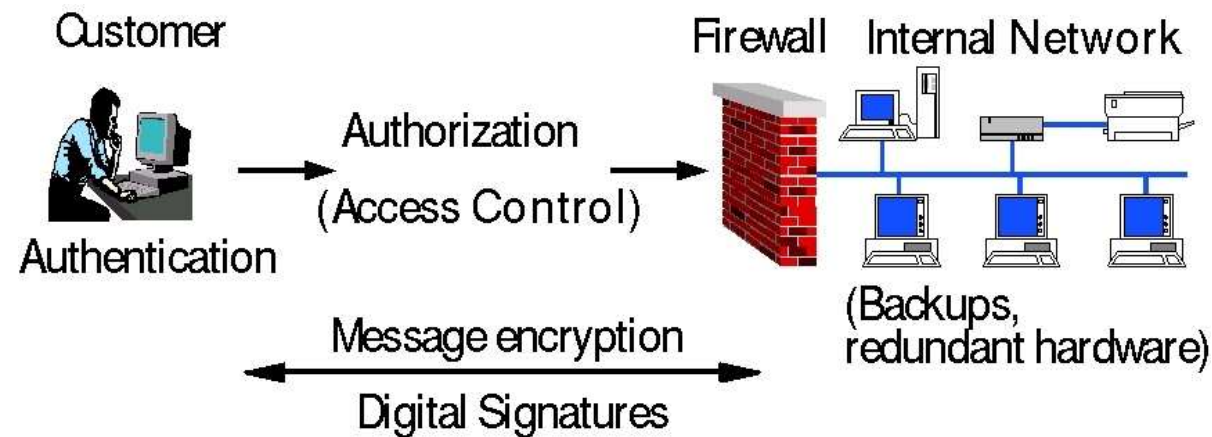
Also called **accountability**

Application specific requirements and combinations, e.g. e-voting

Threats



Security Mechanisms



- Various **mechanisms** are used to achieve **goals**.
- Designing adequate mechanisms is **challenging**.
- One must be cognizant of the **tradeoffs** and **costs** involved.

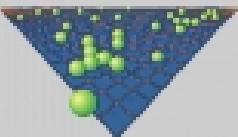
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Security Policies and Models

- A **security policy** defines **what is allowed** (actions, data flow, etc.) typically by a relationship between **subjects** and **objects**.
- A **security model** is a **(+/- formal) description** of a policy and mechanisms, usually in terms of system **state** or sequences of states (**traces**).
- **Security verification** proves wrt. model that **mechanisms enforce policy**
- Models usually focus on **specific characteristics** of the reality (policies).
- Types of (formal) security models
 - **Access Control** models
 - **Information Flow** models
 - **Cryptoprotocol** models



What are Formal Methods?

- A **language** is **formal** if it has a well-defined syntax and semantics.
Examples: Predicate logic, automata, λ -calculus, process algebra, ...

- A **model** is **formal** if it is specified with a formal language.

Example:

$$\forall x. \textit{bird}(x) \rightarrow \textit{flies}(x) \quad \textit{bird}(\textit{tweety})$$

- A **proof** is **formal** if it is done using a deductive system (i.e., a set of precise rules governing each proof step).

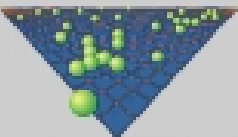
Examples: Tableau calculus, axiomatic calculus, term rewriting, ...

- A formal proof is **machine-assisted** if it is performed, or at least checked, by an IT system.

Examples: OFMC (model checker), Isabelle (theorem prover)

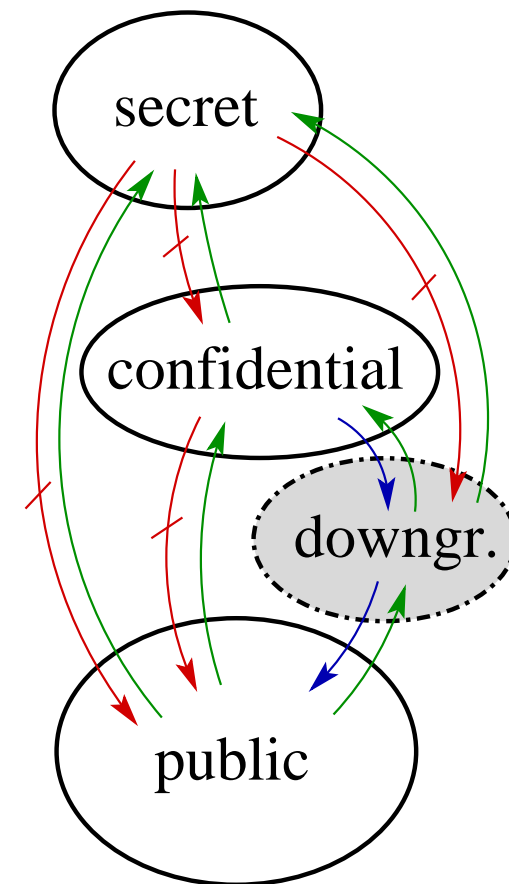
Access Control models

- Discretionary vs. mandatory AC models.
- Various types of models:
 - Models can capture policies for **confidentiality** (Bell-LaPadula) or for **integrity** (Biba, Clark-Wilson).
 - Some models apply to **static policies** (Bell-LaPadula), others consider **dynamic** changes of access rights (Chinese Wall).
 - Security models can be **informal** (Clark-Wilson), **semi-formal**, or **formal** (Bell-LaPadula, Harrison-Ruzzo-Ullman).



Information Flow models

- Identify domains holding information
 - Specify allowed **flow between domains**
 - Check the **observations** that can be made about state and/or actions
 - Consider also **indirect and partial flow**
-
- Classical model: Noninterference (Goguen & Meseguer)
 - Many variants: Non-deducability, Restrictiveness, Non-leakage, ...



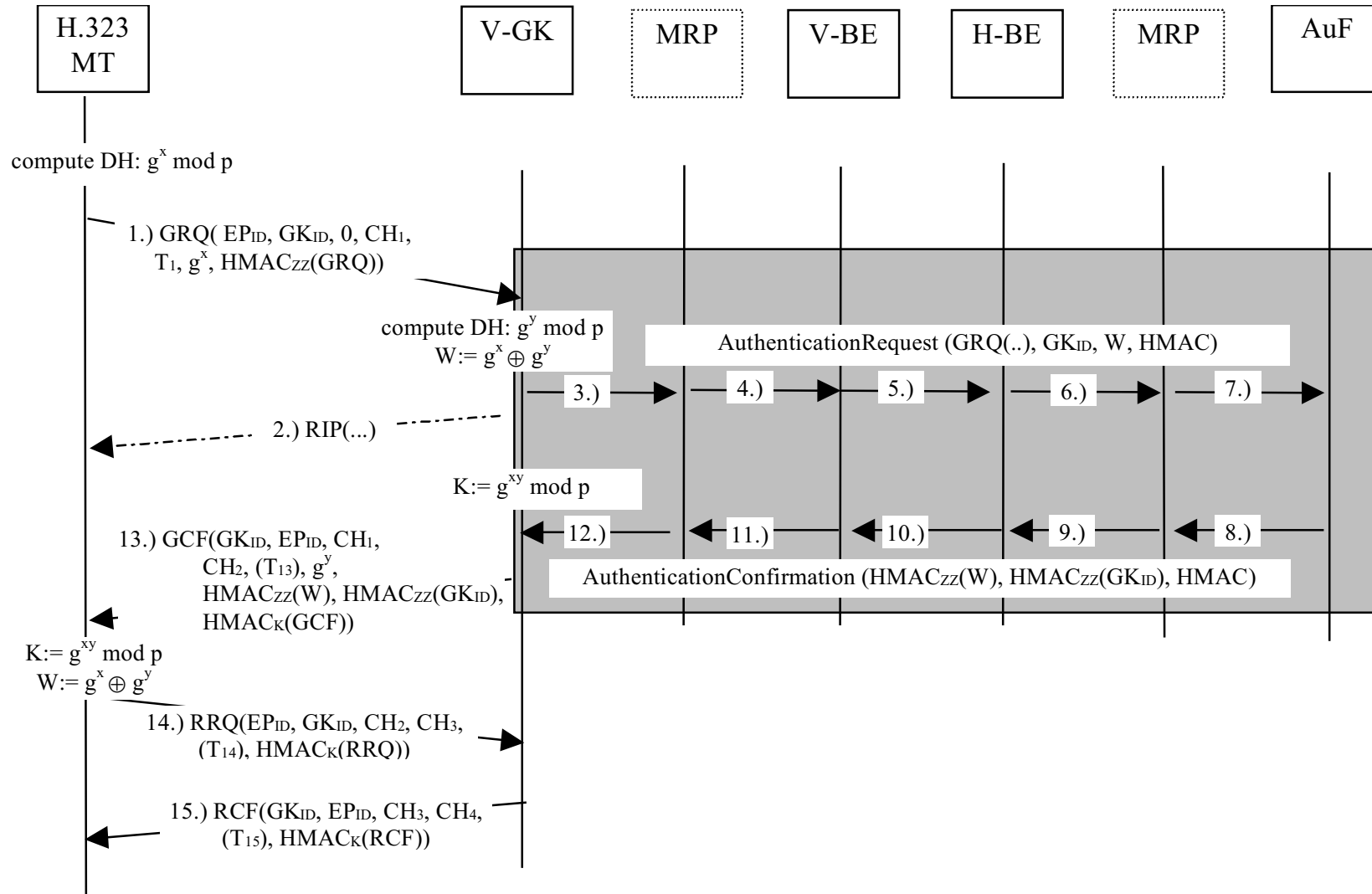
Cryptoprotocol models

- Describe **message traffic** between processes or principals



- Take **cryptographic operations** as **perfect** primitives
- Specified with by domain-specific languages
- Describe **secrecy, authentication, ...** goals
- Are typically verified **automatically** using model-checkers

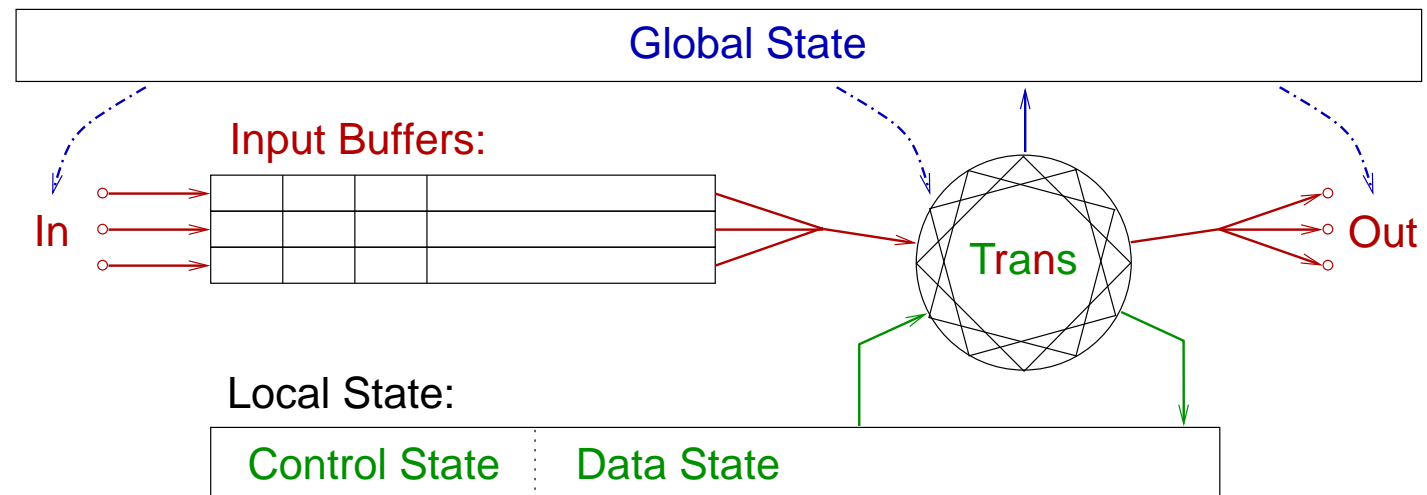
H.530: Authentication for Mobile Roaming



Two vulnerabilities found and corrected. Solution patented.

Interacting State Machines (ISMs)

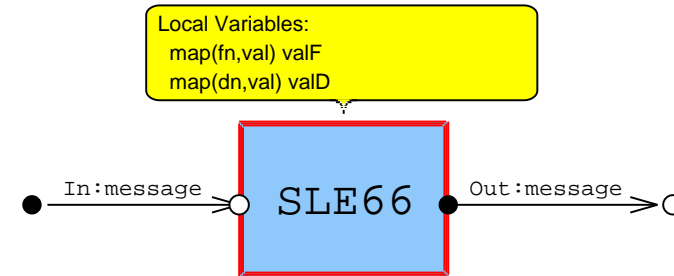
Automata with (nondeterministic) **state transitions** +
buffered i/o simultaneously on multiple connections
 ISM system may depend on **global state**



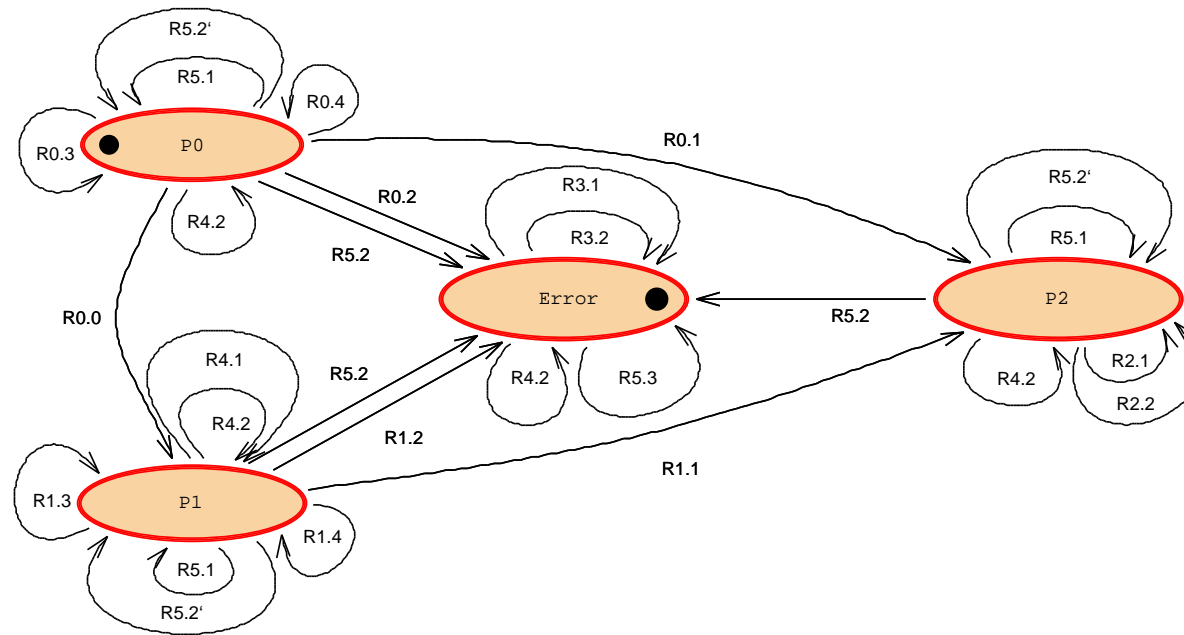
Applicable to a large **variety of reactive systems**

LKW Model of Infineon SLE 66 Smart Card

System Structure Diagram:



State Transition Diagram (abstracted):



First higher-level (EAL5) certification for a smart card processor!

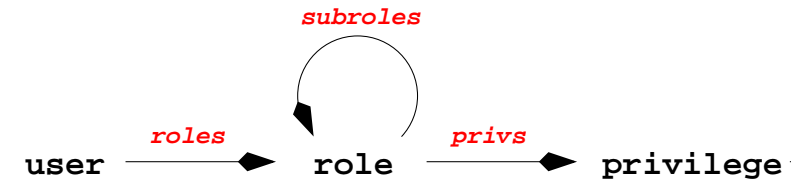
RBAC of Complex Information System

Privileges:

$roles \subseteq user \times role$

$subroles \subseteq role \times role$

$privs \subseteq role \times privilege$



$(u, p) \in roles \circ subroles^* \circ privs$

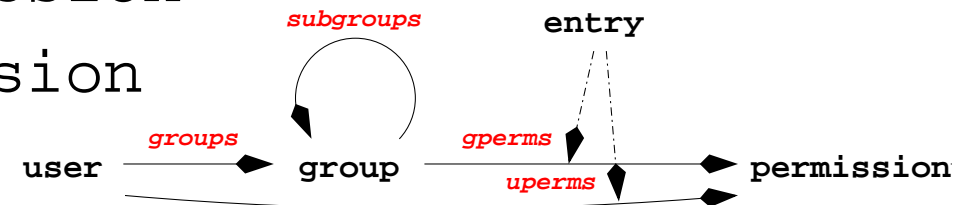
Permissions:

$groups \subseteq user \times group$

$subgroups \subseteq group \times group$

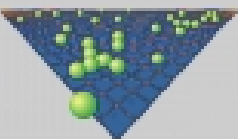
$gperms \subseteq group \times permission$

$uperms \subseteq user \times permission$



$(u, p) \in (groups \circ subgroups^* \circ gperms(e)) \cup uperms(e)$

“nagging questions” \rightsquigarrow clarification improving specification quality



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Shaping a Formal Model

Choice of Formalism: dependent on ...

- application domain, modeler's experience, tool availability, ...
- formalism should be **simple, expressive, flexible, mature**

Formality Level: should be adequate:

- the more formal, the more **precise**,
- but requires deeper mastering of formal methods

Abstraction Level: should be ...

- high enough to achieve **clarity**
- low enough not to lose **important detail**

refinement allows for both high-level and detailed description

Information Necessary

Overview: architecture and components, e.g. authentication services, PKI

Security-related concepts: actors, objects, states, messages, . . .

Threats/security goals/objectives: which attacks shall be countered.

Described in detail such that concrete verification goals can be set up, e.g. integrity: which contents shall only be generated/modified by whom from when to when, or on transit from where to where

Security mechanisms: their relation to goals and how they are applied,

e.g. who signs which contents for what purpose and where checked.

Described precisely but at high level (no implementation details required),

e.g. abstract message contents/format but not concrete syntax

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CC: Goals & General Approach



Common Criteria

...the standard for Information Security

Goal: Gaining **confidence** in the security of a system

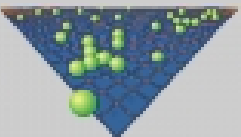
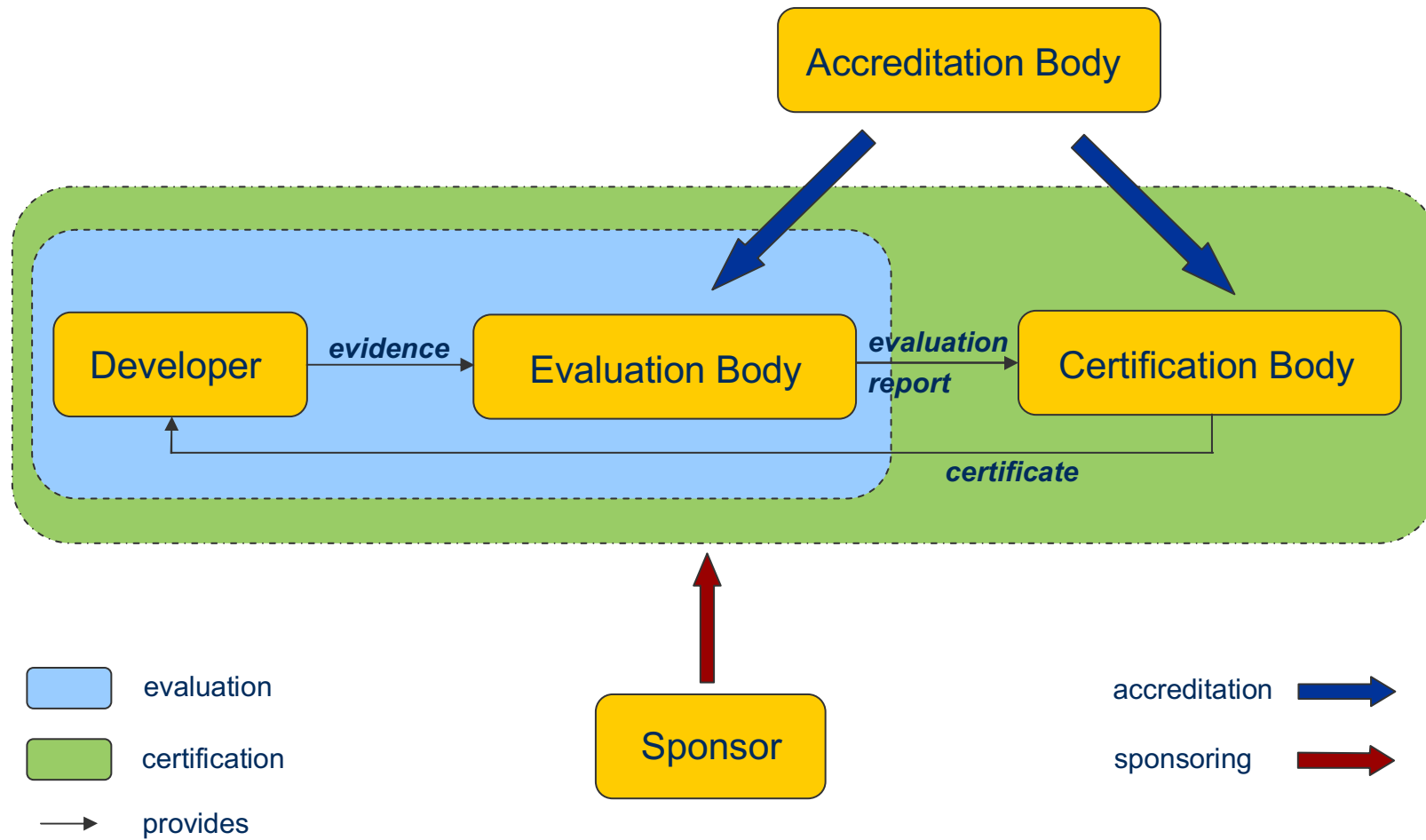
- What are the goals to be achieved?
- Are the measures employed appropriate to achieve the goals?
- Are the measures implemented correctly?

Approach: **assessment** (evaluation) of system security **by neutral experts**

- Understanding how the system's security functionality works
- Gaining evidence that security functionality is correctly implemented
- Gaining evidence that the integrity of the system is kept

Result: Successful evaluation is awarded a **certificate**

CC: Process Scheme



CC: Security Target

- Definition of the **Target of Evaluation (TOE)** and separation from its environment
 - Definition of the TOE's security threats, **objectives** and requirements
 - Introduction of **TOE Security Functions (TSF)**: measures intended to counter the threats
 - Determination of **Evaluation Assurance Level (EAL)**
- ⇒ The Security Target is **the document** to which all subsequent evaluation activities and results refer!
- ⇒ Interpretation of results is only reasonable if referring to the ST context

CC: Evaluation Assurance Levels

EAL1: functionally tested

EAL2: structurally tested

EAL3: methodically tested and checked

EAL4: methodically designed, tested, and reviewed,
including **security policy model**

EAL5: semiformally designed and tested
including **formal** security policy model

EAL6: semiformally verified design and tested

EAL7: formally verified design and tested

Increasing requirements on scope, depth and rigor

CC: EAL example: EAL5

In red: additional requirements compared to EAL4

- **Complete** source code is subject to analysis
- **Formal** security policy model
- **Semiformal description techniques**
- **Modular** design
- Documentation of developer's tests up to **low-level** design
- Vulnerability analysis refers to **moderate** attack potential
- **Covert channel analysis**
- **Comprehensive** configuration management

CC: How to scale an Evaluation

- Separation of TOE and TOE environment
- Detail level of TOE summary specification
- Definition of security objectives
- Definition of security functional requirements
- Strength-of-function claims
- EAL selection



Why are Formal Security Models so useful?

A **formal security model** is an abstract **formal** description of a system (and its environment) that focuses on the relevant security issues.

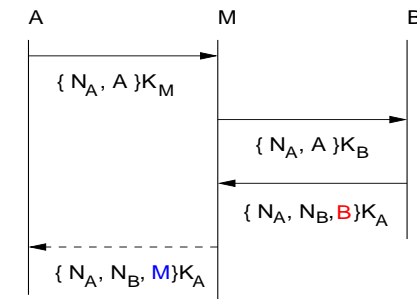


Abstraction

→

Interpretation

←



- **improves understanding of security issues** by
 - abstraction: concentration on the essentials helps to keep overview
 - systematic approach: generic patterns simplify the analysis
- **prevents ambiguities, incompleteness, and inconsistencies** and thus enhances quality of specifications
- **provides basis for systematic testing or even formal verification** and thus validates correctness of implementations